Cooperative Communications and Networking

**Chapter 4** 

#### **Relay Channels and Protocols**

# Outline

- Cooperative communications: a paradigm shift
- Definitions and system model
- Relaying strategies
  - Amplify & forward
  - Selection decode & forward
  - Incremental relaying
- Hierarchical cooperation

# A Paradigm Shift

- Future generations of wireless systems promise very high spectral efficiencies
- MIMO is a key enabler for achieving higher data rates
- The number of antennas on any device is limited by the device size
- Cooperation offers a new communication paradigm in which the number of antennas virtually increases with the number of terminals
- Such a virtual antenna system can offer significant capacity gains

### **Relaying Strategies**



### **Definitions and System Model**

- A channel is:
  - Fast fading if each packet encounters several channel realizations
  - Slow fading if each packet encounters one channel realization
- Quasi-static Rayleigh fading channel is assumed
- Channel state information (CSI) at receiver
- Gaussian channel input, unless otherwise stated
- Half-duplex devices

#### Definitions

• The multiplexing gain

$$r = \lim_{SNR \to \infty} \frac{R(SNR)}{\log(SNR)}$$

• The diversity gain

$$d(r) = -\lim_{SNR \to \infty} \frac{\log p_{out}(r \log SNR)}{\log SNR}$$

#### Definitions (Diversity and Multiplexing Gains)

• The spatial multiplexing gain of a  $n_t \times n_r$  MIMO scales like

 $\min(n_t, n_r) \log(SNR) \longrightarrow For 0 \text{ div gain}$ 

• The diversity gain of  $n_t \times n_r$  MIMO is

$$SNR^{-n_t n_r}$$
 — For 0 multiplexing gain

There is a Tradeoff

#### **Relevant Performance Metrics**

### Outage probability

 Probability of not meeting specified transmission rate R due to poor channel realizations

 $P_{out} = \Pr[I(x, y) < R]$ 

 I(x, y) is the maximum mutual information between input and output signals

# Diversity gain

 Reduction in outage probability due to SNR improvement

$$\nabla = -\lim_{\gamma \to \infty} \frac{\log P_{out}}{\log \gamma} \qquad \qquad \gamma : \text{signal-to-noise ratio}$$



- input x, noise terms n<sub>sd</sub>, n<sub>sr</sub>, n<sub>rd</sub>
- Rayleigh fading channels
  - Nodes know channel coefficients

 $h_{sd} \sim R(\sigma_{sd})$  $h_{sr} \sim R(\sigma_{sr})$  $h_{rd} \sim R(\sigma_{rd})$ 

# Fixed Relaying: Amplify and Forward <sub>R</sub> SO Phase Amplify y<sub>sr</sub> and forward to destination $y_{rd} = h_{rd}q(y_{sr}) + n_{rd}$ $=h_{rd}\left|\frac{\sqrt{P}}{\sqrt{P|h_{sr}|^2+N_0}}\left(\sqrt{P}h_{sr}x+n_{sr}\right)\right|+n_{rd}$ $= \left(\frac{\sqrt{P}}{\sqrt{P|h_{sr}|^2 + N_0}}\sqrt{P}h_{rd}h_{sr}\right)x + \left(\frac{\sqrt{P}}{\sqrt{P|h_{sr}|^2 + N_0}}h_{rd}n_{sr} + n_{rd}\right)$ Both y<sub>sd</sub> and y<sub>rd</sub> are fused with maximum ratio combining at destination to maximize SNR $y = \left[\frac{\sqrt{P}h_{sd}^{*}}{N_{0}}\right]y_{sd} + \left|\frac{\frac{\sqrt{P}}{\sqrt{P|h_{sr}|^{2} + N_{0}}}\sqrt{P}h_{rd}^{*}h_{sr}^{*}}{\left(\frac{P|h_{rd}|^{2}}{P|h_{sr}|^{2} + N} + 1\right)N_{0}}\right|y_{rd}$

### Fixed Relaying: AF

Instantaneous SNR at destination

$$\gamma = \gamma_{sd} + \gamma_{rd} = \frac{P|h_{sd}|^2}{N_0} + \frac{1}{N_0} \frac{P^2|h_{rd}|^2|h_{sr}|^2}{P|h_{sr}|^2 + P|h_{rd}|^2 + N_0}$$

Mutual information

#### Source and relay each take half of resource

$$I_{AF} = \frac{1}{2} \log \left( 1 + \gamma_{sd} + \gamma_{rd} \right) = \frac{1}{2} \log \left( 1 + \Gamma \left| h_{sd} \right|^2 + f \left( \Gamma \left| h_{sr} \right|^2, \Gamma \left| h_{rd} \right|^2 \right) \right) \qquad \Gamma = P / N_0$$
  
Outage probability  
$$f(a,b) = \frac{ab}{a+b+1}$$

Outage probability

$$P_{out} = \Pr[I_{AF} < R] \approx \left(\frac{\sigma_{sr}^2 + \sigma_{rd}^2}{2\sigma_{sd}^2(\sigma_{rd}^2\sigma_{sr}^2)}\right) \left(\frac{2^{2R} - 1}{\Gamma}\right)^2 \qquad \nabla = 2$$

#### Fixed Relaying: AF



### Fixed Relaying: Decode and Forward

- Decode, re-encode, then forward signal
  - Re-encoded signal might be incorrect



Mutual information limited by the weakest link between the source-relay and the combined source-destination and relay-destination.

$$I_{DF} = \frac{1}{2} \min \left\{ \log \left( 1 + \Gamma |h_{sr}|^2 \right), \log \left( 1 + \Gamma |h_{sd}|^2 + \Gamma |h_{rd}|^2 \right) \right\}$$

Outage probability

$$P_{out} = \Pr[I_{DF} < R] \approx \frac{1}{\sigma_{sr}^2} \frac{2^{2R} - 1}{\Gamma}$$

 $\nabla = 1$ 

### Adaptive Relaying: Selective DF

DF if source-relay link exceeds SNR threshold



Outage probability

$$P_{out} = \Pr[I_{SDF} < R] \approx \left(\frac{\sigma_{sr}^2 + \sigma_{rd}^2}{2\sigma_{sd}^2(\sigma_{rd}^2 \sigma_{sr}^2)}\right) \left(\frac{2^{2R} - 1}{\Gamma}\right)^2 \qquad \nabla = 2$$

### Adaptive Relaying: Incremental

Feedback from the destination if source signal is transmitted successfully



Overall transmission rate

$$\overline{R} = R \Pr\left[\left|h_{sd}\right|^2 \ge \frac{2^R - 1}{\Gamma}\right] + \frac{R}{2} \Pr\left[\left|h_{sd}\right|^2 < \frac{2^R - 1}{\Gamma}\right] \ge \frac{R}{2}$$

• Outage probability  

$$\Pr[I_{IR} < R] = \Pr[I_{Direct} < R]\Pr\left[I_{AF} < \frac{R}{2} \mid I_{Direct} < R\right]$$

$$\approx \left(\frac{\sigma_{sr}^2 + \sigma_{rd}^2}{\sigma_{sd}^2 \left(\sigma_{rd}^2 \sigma_{sr}^2\right)}\right) \left(\frac{2^{\overline{R}} - 1}{\Gamma}\right)^2 \qquad \nabla = 2$$

#### Outage Versus SNR for fixed Rate



#### **Outage Versus Rate for fixed SNR**



### **Hierarchical Cooperation**

- Motivation: n nodes within communication range want to communicate at rate R(n) per node
  - Each node is assumed to have a specific and fixed receiver
- Problem: Long-range transmission using large power causes too much interference
  - Aggregate throughput T(n) = nR(n) degrades as n grows

### **Hierarchical Cooperation**

- What is the maximum aggregate throughput that can be achieved in a wireless network?
- In a n-node network, nearest neighbor multihop transmission has a network capacity that is upper bounded as [Gupta/Kumar, 00]

$$T(n) \le O(\sqrt{n})$$

• Throughput per node drops as the density of the nodes increases as

$$R(n) = \frac{T(n)}{n} \le O(\frac{1}{\sqrt{n}})$$

• What can cooperation help?

### **Hierarchical Cooperation**

• *Hierarchical cooperation* can achieve linear scaling!

$$nR(n) \approx O(n)$$

• This means: As the network density increases, the throughput/node does not degrade!

• Network throughput is bounded above by

 $nR(n) \le O(n\log n)$ 

# Hierarchical Cooperation (2)

- Hierarchical cooperation scheme can achieve linear scaling of network capacity in ad hoc networks if it leverages [Özgür/Lévêque/Tse, 07]
  - Clustering (spatial reuse)
  - Long-range MIMO transmission across clusters (high spatial multiplexing gain)
- Hierarchical cooperation scheme recursivel applies 3-phase protocol each iteration, which are
  - Phase I: Setting up transmit cooperation
  - Phase II: MIMO transmission
  - Phase III: Cooperate to decode



#### Phase 1: Setting Up Transmit Cooperation

- The network is divided into clusters each with M nodes
- Clusters operate in parallel following the 9-TDMA scheme
- Each node as a source divides a block of length LM bits into M sub-blocks, and disseminates these sub-blocks among the nodes inside its clusters

#### Phase 1: Setting Up Transmit Cooperation



#### Phase 1: Local transmit cooperation in clusters



#### Phase 2: MIMO long range transmission between clusters



#### Phase 3: Local receive cooperation in clusters



# Why Does it Work?

- Via clustering one can achieve spatial reuse
- This is because clusters that are well separated can work in parallel
- The nodes operating in each cluster can lower their power accordingly to limit the interference to other clusters working in parallel
- Via long-range MIMO we can achieve high spatial multiplexing gain

### Analysis: Does It Work?

• Key Lemma: If there exists a scheme that can achieve an aggregate throughput of

 $T(n) \ge K_1 n^b, \qquad 0 \le b < 1$ 

with upper-bounded average power, then one can construct another scheme that achieves 1

$$T(n) \ge K_2 n^{\overline{2-b}}$$

with the same average power

#### Analysis: Does It Work?



• Performance is improved, particularly for small b

#### Proof Sketch of Key Lemma: Analysis of Phase 1

- Each node in a cluster has LM bits to transmit
- Each node scales its power proportional to the cluster size
   Interference at each node is bounded
  - Aggregate throughput of  $K_1 M^b$  can be achieved
- Total time slots to complete Phase 1 is therefore

$$T_{1} = \underbrace{2M \times \frac{LM}{K_{1}M^{b}}}_{\text{required time slots for}} \times 9 = \frac{18L}{K_{1}}M^{2-b}$$
  
required time slots for  $\therefore$  9-TDMA Scheme  
total traffic  $\Box M \times LM$ 

(independent of n due to parallel processing)

### **Analysis of Phase 2**

- Long-range MIMO transmission takes place
- An MxM multiple antenna MIMO system can transmit O(M) bits in 1 time slot
  - Essential step to achieve high capacity scaling
- Time slots required to complete all MIMO transmission is proportional to n

$$T_2 = 2Cn$$

### **Analysis of Phase 3**

- Each node delivers the observation it receives from Phase 2 to the destination node
- Similar to Phase 1, clusters work in parallel using 9-TDMA, and the required number of time slots is given by

$$T_{3} = \frac{18CQ}{K_{1}} M^{2-b}$$

#### Combining Phase 1, 2, and 3

• Total number of time slots is given by

$$T_{t} = T_{1} + T_{2} + T_{3} = \frac{18L}{K_{1}}M^{2-b} + 2Cn + \frac{18CQ}{K_{1}}M^{2-b}$$

Aggregate throughput will be thus as given in the lemma

$$T(n) = \frac{nML}{T_t}$$

$$= \frac{nML}{\frac{18L}{K_1}M^{2-b} + 2Cn + \frac{18CQ}{K_1}M^{2-b}} \qquad \text{Tradeoff between cooperation overhead (T1+T3) and MIMO performance}$$

$$\geq K_2 n^{\frac{1}{2-b}}, \text{ for some } K_2 \text{ if } M = n^{\frac{1}{2-b}} \text{ is chosen}$$

(T1+T3) and MIMO

### **Achieving Linear Capacity Scaling**

- Idea: recursively apply the 3-phase protocol in Phase 1 and 3 to achieve higher throughput (hence "Hierarchical")
  - Recursion starts with smaller areas and grows until including the whole network
- Theorem: For any  $\varepsilon > 0$  there exists a constant  $K_{\varepsilon}$  such that, with high probability, an aggregate throughput of  $T(n) \ge K_{\varepsilon} n^{1-\varepsilon}$  is achievable
- Eventually, we get the linear capacity scaling! This means that cooperative communication can help to make the network scalable.

### **Achieving Linear Capacity Scaling**

• Hierarchical Cooperation



PHASE 1			PHASE 2	PHASE 3		
PHASE 1	PHASE 2	PHASE 3		PHASE 1	PHASE 2	PHASE 3